

DESIGN AND PERFORMANCE ANALYSIS OF 32-BIT ADDERS FOR AREA EFFICIENT VLSI APPLICATIONS

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ABSTRACT:

This paper represents the basic process like firstly addition subtraction can be done using various types of binary adders with dissimilar addition times, area consumption in any digital processing application. Parallel prefix adders are considered to be one of the fastest adders that had been designed and developed. Parallel prefix adders were established as the most efficient circuits for binary addition. This paper investigates the performance of three different parallel prefix adders namely Kogge Stone Adder(KSA), Brent Kung Adder(BKA) and Ladner Fischer Adder(LFA). In this paper the key contribution is the information about the structure of the parallel prefix adders. All the simulations and synthesis results can be noted using Xilinx ISE 14.2i tool.

Keywords: Parallel Prefix Adders, Kogge Stone Adder(KSA), Ladner Fischer Adder(LFA), Brent Kung Adder(BKA), Carry look ahead adder, carry generation stage.

I.INTRODUCTION

Generally, the basic process such as addition, subtraction, division that can be done by using different types of binary adders in any digital based processors and control systems. The adder performance is used in the device is only measured the high speed and accuracy of the system. Previously the processors are used 32-bit carry adders like Ripple carry Adder(RCA), Carry Propagate Adder(CPA), and Carry Look ahead Adder(CLA) with different delay and area consumption.

How fast the carry reaches for every single bit position, from which the delay of any binary adder is calculated. But the above existing 32-bit basic carry adders having a delay value in higher order bits because each level of adder has to wait for the previous carry result. Due to the above problem of 32-bit basics existing carry adders, in today's world of technology, PPA is well suitable designed adder for high speed addition process with less delay in VLSI technology.

PPA is one of the most popular designs and provides good negotiation amongst area and delay. The low order PPA is designed at earlier like 8-bit and 16-bit. This paper consisting of three additions first addition: describes briefly about PPA, in second addition: explain the design of 32-bit proposed parallel prefix(Kogge Stone Adder, Brent Kung Adder, Ladner Fischer Adder), and in third addition it express the simulation results(waveform and reports) of 32-bit PPA which we have designed in the previous section with the performance aspects(delay and area). The last section is concluded that Kogge Stone Adder performance is the best among the other adders with low area and less delay from the analyzed results of 32-bit PPA

1. Parallel Prefix Adder

Now a days, to avoid the higher delay problem of existing carry adders the PPA is used which is simply the modified design form of CLA. The Prefix adders can be designed in many different ways based on the different requirements and the production of carries[6]. Recently, use the tree structure form of adder to store the speed of addition function in any kind of processors. PPA are fastest adders with tree structure based and used for high performance arithmetic processes in successive industries and DSP laboratories [7].

The PPA's are also called as logarithmic delay adders because the delay value is established using logarithmic functions [8]. Addition in PPA can be processed using three main actions such as Pre-computation (P and G signal generation), Prefix-computation (carry signals group generation), Post-computation (Sum signal generation) [9].

1.1. Pre-Computation

In the pre-processing stage, propagate functions and generate functions are calculated depends up on the given input signals[10].The propagate functions are carried out by the equation (1).

$$P[j]=X[j] \oplus Y[j] \tag{1}$$

It is stated that X and Y are the input signals that composed by XOR logic gate. The generate functions are carried out by the equation(2).

$$G[j]=X[j] \cdot Y[j] \tag{2}$$

It is stated that X and Y are the input signals that composed by AND logic gate. Since the above equations (1) and (2) are done in parallel, it does not increase a significant calculation of area consumption and delay fully depends upon the bit size which is desired at the input [11].

1.2. Prefix- Computation

This prefix computation stage, calculates the carry signal groups directly, which uses the input and values which measured from the first stage. Carry signals generation uses the more than two inputs for which the delay is automatically increased in this process[12].The carry propagation function and carry generation function [13] is measured by the equations (3) and (4).

$$\text{Carry } P = P [j] \cdot P [j + 1] \tag{3}$$

$$\text{Carry } G = (G [j] \cdot P [j + 1]) + G [j + 1] \tag{4}$$

1.3. Post-Computation

In this stage, the sum result is generated by an Ex-OR operation that uses the values of carry generation stage (prefix-computation). The last sum operation is calculated by the equation (5)

$$S[j+1]=P[j+1] \oplus C[j] \tag{5}$$

Where C is the last carry signal and P is the propagate function [14].

2. Design and analysis of proposed 32-bit various PPA

To overcome the high delay problem of existing carry adders, this work proposed the design of 32-bit various PPA for less delay and low power VLSI applications. This proposed system consists of two modules: The first module is to design of 32-bit PPA like KSA, BKA, and LFA. The second module is to analyze the performance comparison of PPA on the basis of area, delay, and power.

In this section, we analyze the different technologies of adders to design in the form of parallel prefix, apart from the RCA topology, such as Kogge-Stone, Brent-Kung, and Ladner-Fisher PPA. The important aim is to examine the trade-off between area consumption, delay, and power consumption in the particular PPA depends upon the design performance. All the designs are using a power gating technique to reduce the power consumption [15].

2.1. Kogge Stone Adder (KSA)

Normally, the KSA attains the key role with fast addition operation and it reflects like the prefix form of Carry Look ahead Adder (CLA). Also, this type of PPA entirely decreases the delay time in design to generate the carry signals [16]. Henceforth, this KSA is popularly used in DSP (Digital Signal Processing) laboratory and Control system industries for fast arithmetic function.

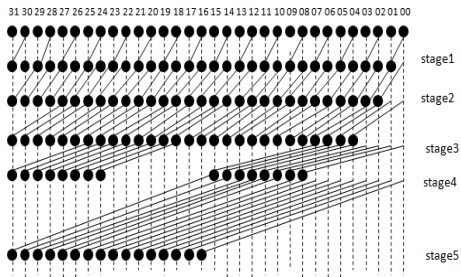


Figure 1: 32-bit Kogge Stone Adder design

The structure of 32-bit KSA design is exposed in Figure 1. This design can be divided into 5 stages [17]. The calculation of Propagate and Generate signals using full adders with carry input that process included in first and second stage. The generation of carry signals which used the values of Propagate and Generate that process included in third and fourth stage. The calculation of sum bits based on the P and carry generation values that is included in the fifth stage [18]. This 32-bit design of KSA is coded by VHDL and viewed the test bench waveform and analyzed the performance and noted the results.

2.2. Brent Kung Adder (BKA)

The BKA calculates the prefixes based on the bit groups.

Initially calculate the prefix values for 2 bit groups. These 2 bit prefix values are used to find the prefix values for the 4 bit groups, that are used to calculate the prefix values for 8 bit groups and so on [19]. Then these prefix values are used to measure the carry out of the particular bit stage. These carries will be used along with the Group Propagate of the next stage to calculate the Sum bit of that stage [20]. Brent Kung Tree will be using $(2 \log_2 N - 1)$ stages for any bit design.

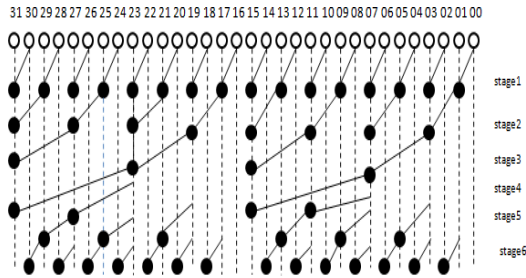


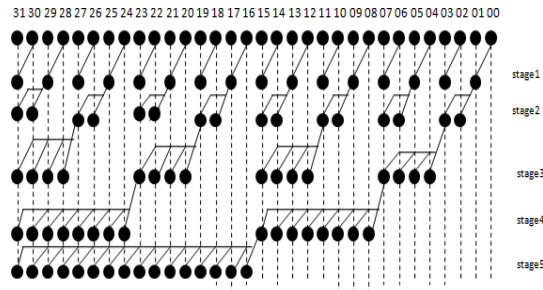
Figure 2: 32-bit Brent Kung Adder design

The structure of 32-bit BKA design is given in Figure 2.

Hence the designing of 32-bit adder takes the number of stages will be 9. The fan-out for every bit stage is limited to 2. The above diagram shows the fan out being reduced and the loading on the advance stages being reduced [21]. This 32-bit design of BKA is coded by VHDL and viewed the test bench waveform and analysed the performance and noted the results.

2.3. Ladner Fischer Adder (LFA)

The LFA tree structures are a family of tree networks between Brent Kung and Sklansky tree. It is very close like to Sklansky PPA, but it calculates the prefix values for odd number bits after that again uses another stage which ripple into



the even locations [22]. At higher order bits, to get improved in speeds the cells must still be properly sized or grouped.

Figure 3: 32-bit Ladner Fischer Adder design

The structure of 32-bit LFA design is exposed in Figure 3. The Ladner Fischer adder is used for high performance arithmetic operation with complicate designs. The LFA consists of black cells and gray cells with 5 stages for 32-bit design. Each black cell encloses only one OR logic gate and two AND logic gates. Each gray cell contains only one AND logic gate [23]. This 32-bit design of LFA is coded by VHDL and viewed the test bench waveform and analysed the performance and noted the results.

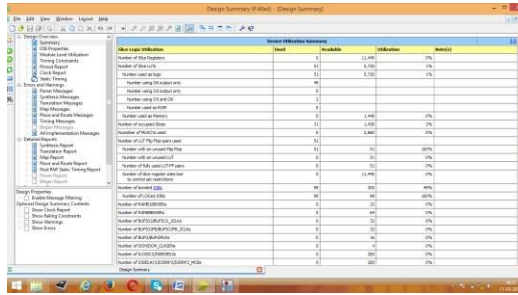


Figure10:AreaConsumptionreportforLFA

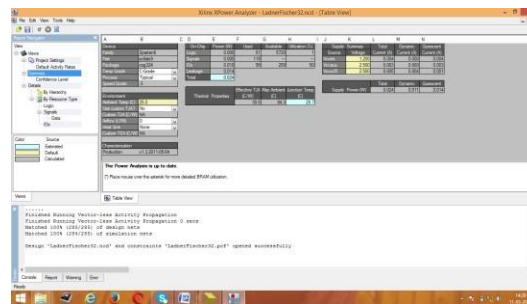


Figure11:PowerAnalyzerreportforKSA

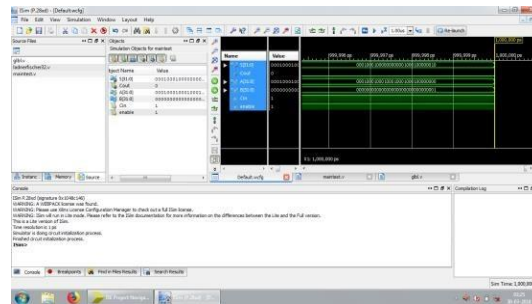


Figure12:Testbenchwaveformresultsof LFA

Area consumption, power analysis, and the test bench waveform of KSA are shown in figures 4, 5, and 6 respectively. Area consumption, power analysis, and the test bench waveform of BKA are shown in figures 7, 8, and 9 respectively. Also the area consumption, power analysis, and the test bench waveform of LFA are shown in figures 10, 11, and 12 respectively.

From the above figures, the comparison results of all three PPA on the aspects of area, delay and power is given in table I. From the analysis, LFA is better due to the less area consumption but the power utilization is more compared to other adders. Normally PPA's have less delay in any processors while doing addition.

Accordingancewithlowpowerapplication,KSAismoresuitabletoless power utilization in any digital based processors.

Table 1: Comparisonresults of area, delay and power for varioustypes of PPA

AdderTypes	LUT's Used	Delay(ns)	Power (mW)
KoggeStoneAdder	189	6.483	21
BrentkungAdder	156	6.489	24
LadnerFischerAdder	51	21.879	24

Conclusion

In this paper, an efficient 32 bit Parallel Prefix adders like KSA, BKA, LFA is designed. This proposed 32 bit adder addition operation offers a great advantage in reducing delay. For low power VLSI applications, also the designed adders are compared on the basis of power, area consumption, and delay. The synthesis results reveal that among the proposed adders, KSA is achieved some saving of power-delay product due to less power utilization. But the area delay product is little increased, compared to other adders due to high area consumption. For decreasing the complexity at all performance aspects, further optimization techniques can be achieved on the performance parameters that will be the future work of the paper.

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