

SOFTWARE DEFINED TESTBED USING MININET EMULATOR

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ABSTRACT

Software Defined Networking (SDN) is changing the way future networks will function. Already Big companies like Google, facebook are using SDN. SDN is also the emerging research area. For doing research we need SDN testbed such as Geni, Emulab etc or very costly sdn switches. Since most of the time, getting resources at testbeds or purchasing hardware for SDN is difficult, we need some other solution. For performing research, Mininet emulator has been created. Mininet can be easily installed on Desktop, Laptop, Physical server or as a virtual machine thus enabling you to perform simple as well as complex experiments on limited resources available. Mininet is an emulator which uses light weight virtualization technology to create simple as well as complex networks. It enables you to create a network consisting of few hosts and switches as well as thousands of hosts and switches by using simple commands. Other solutions for doing research could be using simulators such as NS3 but the problems with simulator is that you cannot run real code on them whereas an emulator allows you to run real, unmodified code. Other choice is creating hardware test beds which can prove very costly and time consuming. Mininet offers easiness of use, good performance, scalability, accuracy. In this paper, we will paper, first section will be introduction to Mininet, Second section about SDN Architecture, 3rd Section will discuss about different mininet topologies, 4th section about using mininet with network, Fifth section will deal with using Mininet advantages & disadvantages & last section will be conclusion.

Keywords: Mininet, OpenFlow, Software Defined Networking, POX Controller

I. INTRODUCTION

Mininet is an emulator which allows you to create various topologies consisting of hosts, switches, controllers and links between them. Mininet [1] is very easy to install, use and very light weight in resource usage. It can be used as SDN test bed for performing SDN [2] experiments. It is very powerful tool in testing SDN networking applications. It is written in python and uses light weight virtualization for creating and isolating network resources.

Mininet offers many features that are available in other simulators, emulators and testbeds. There are many hardware testbeds such as VINI [3], FIRE [4], Emulab [5], and GENI [6] available. Using these testbeds is difficult and if you want to create testbeds of these types then it is very costly and time consuming process. There are simulators such as NS3 [7] and Estinet [8] available, but you cannot run real code on these simulators. On the other hand, Mininet is readily available, very inexpensive, quickly configurable, runs unmodified code.

You can use mininet to create pre defined topologies or custom topologies. Another advantage is that it is open source.

II. SOFTWARE DEFINED NETWORKING ARCHITECTURE

SDN is going to change the ways future networks are going to get implemented. All the network devices consist of control plane and data plane. Control plane contains the control logic and data plane performs the actual forwarding of the packets. But the problem with these devices is that they are tightly integrated. You cannot buy them separately. Their functionality is rigid. These devices are classified as hub, router, switches, load balancers, proxy servers, firewall, intrusion detection, intrusion prevention systems. The devices are vendor specific and very costly. There is very large time gap between adding new features in these devices. Also it depends on the will of vendors whether they are going to add the feature or not and most importantly when. Most of the time, their decision is based on financial considerations. SDN has decoupled these planes. Now you can acquire & use data plane and control plane separately. The data plane is available in both the physical and virtual form. The control plane is also called controller (network OS). Various controllers such as NOX [9], POX [10], Beacon [11], Floodlight [12], OpenDayLight [13], and RYU [14] are available. All of these controllers are open source and free. The difference is mainly in which languages they are written and what openflow version they support. Openflow is protocol for communication between control plane and data plane. POX, RYU are written in python while Floodlight, OpenDaylight in Java. While Pox supports Openflow version 1.0 while RYU supports versions from 1.0 to 1.5.

The architecture of SDN consists of 3 components, Openflow Controllers, switches and Openflow protocol as shown in Fig.1. Southbound openflow [15] protocol is used for control plane and data plane communication. The connection between both is maintained through secure channel. Openflow based devices consists of flow table. Flow tables contain flow entries having matching fields, actions and flow statistics. The match fields are based on layer 1 to layer 4 fields such as source mac, destination ip, port etc. Actions could be forward, drop. Packets received at switch are matched against flow entries. If match is found, action is performed according to flow entry. If no match is found, then packet is send to the controller where the application logic will decide what to do with the packets. Then controller will add entries into flow table of switch. It is the controller's responsibility to manage entries in the switch table.

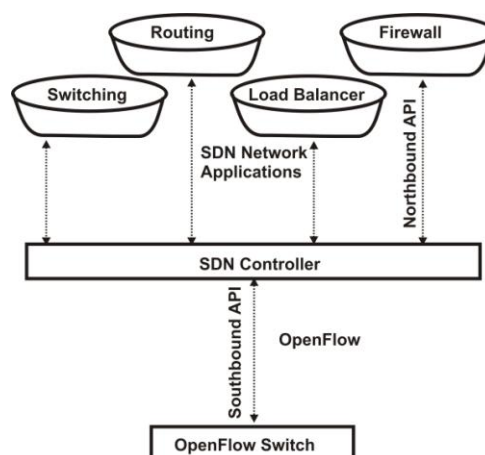


Fig. 1. SDN architecture

III. MININET TOPOLOGIES

There are many default topologies available in mininet such as minimal, linear, tree, single. Here we will be discussing these different topologies. Topology in mininet consists of hosts, switches, controller and links between them as shown in Fig. 2. Understanding naming conventions is required. Host are named from h1 to hN, switches are from s1 to sN. Names of the interfaces in hosts are created by prefixing host name followed by Ethernet name beginning with 0. Host h1's first interface will be h1-eth0, second h1-eth1, and third h1-eth2 and

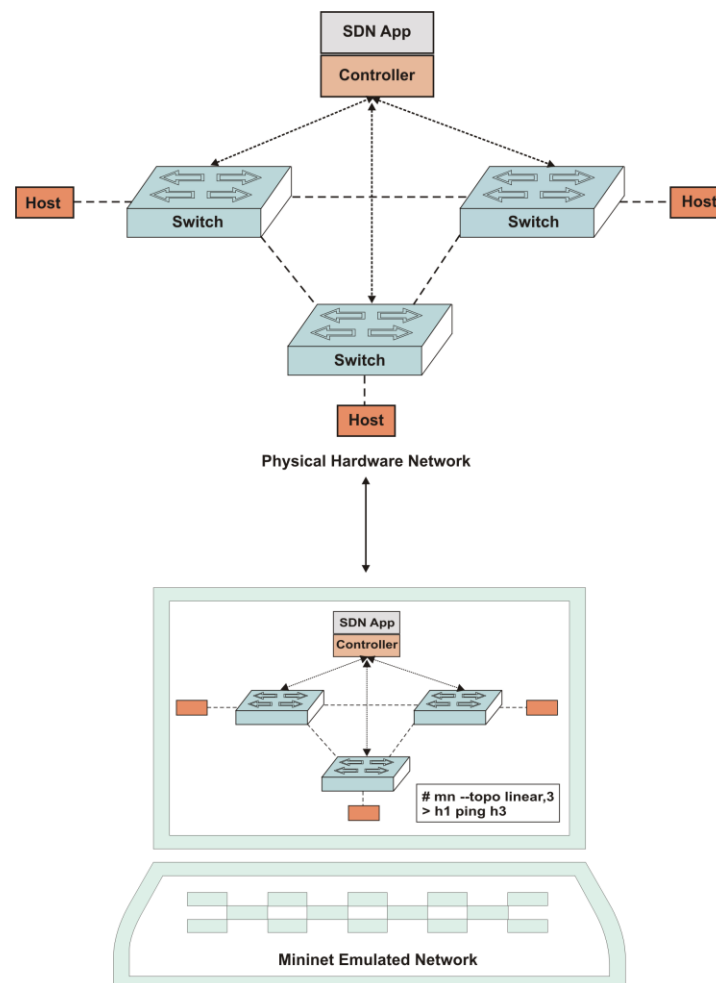


Fig. 2 mininet emulator

so on. In switches, the number starts from 1. So Switch ports are named s1-eth1, s1-eth2, and s1-eth3 and so on.

3.1 Mininet Topologies

3.1.1 Minimal Topology

mn --topo minimal

It will create a topology consisting of 1 switch s1 and 2 hosts h1 and h2. The hosts will be having default ip address of 10.0.0.1 and 10.0.0.2 as shown in Fig. 3.

3.1.2 Single Topology

mn --topo single,3

It creates 1 switch and n hosts. It creates links between switch and n hosts as shown in Fig. 4

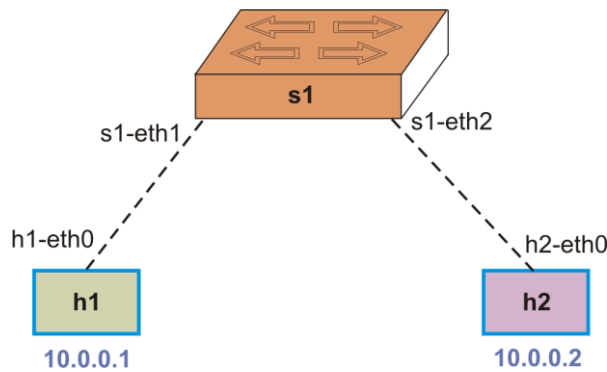


Fig. 3 minimal topology

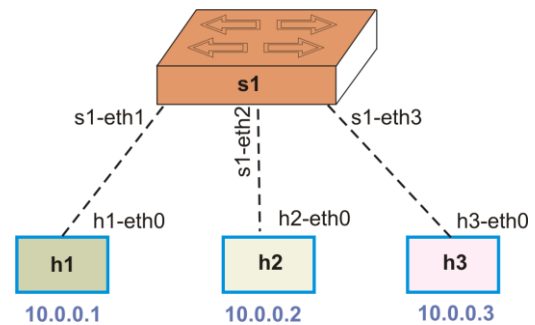


Fig. 4 single topology

3.1.3 Linear Topology

mn --topo linear,3

It consists of n switches and n hosts. It creates link between hosts, switch & between switches as shown in Fig.

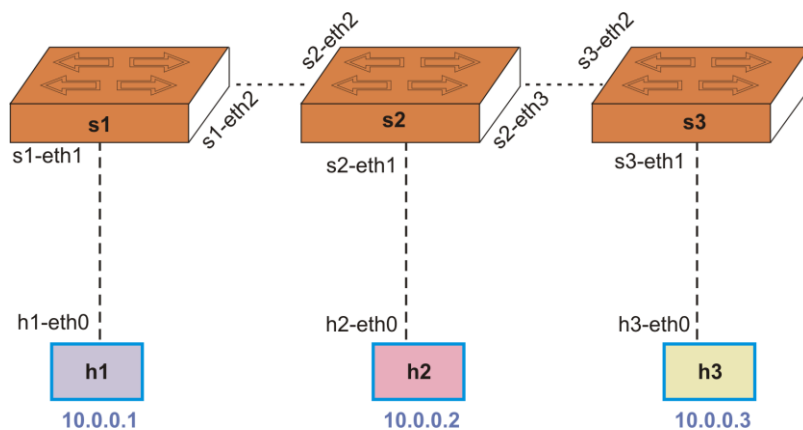


Fig. 5 linear topology

5.

3.1.4 Tree Topology

mn --topo tree,2

It creates topology having n levels and default 2 hosts are attached with each switch as shown in Fig. 6

3.1.5 Custom Topology

We can create custom topologies by writing few line of python code. We can create any type of complex topology.

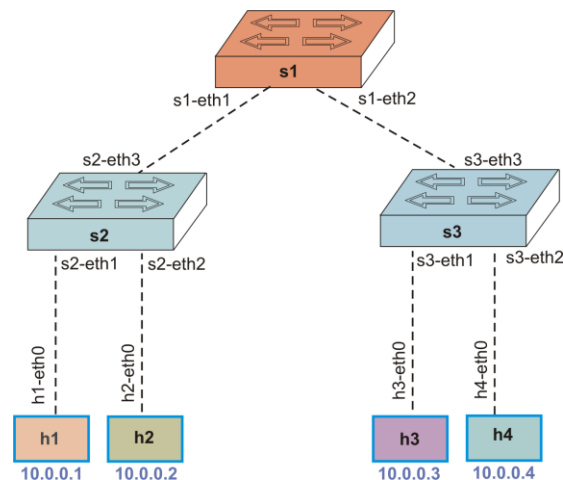


Fig. 6 tree topology

IV. USING MININET WITH NETWORK APPLICATIONS

We can create networking applications such as hub, switch and load balancer using controllers. We can use mininet to test these applications. We are here using pox controller. Pox controller is written in python. By writing few lines of python code, we can write these applications. The mininet virtual machine that we downloaded from mininet web site already contains hub, switch and load balancing applications. These are very simple applications. We can create very complex applications like firewall, intrusion detection system, intrusion prevention systems and then test these applications on mininet. Whenever you are creating applications, we should be aware of the logic behind these devices i.e. we must know how these devices work, only then we can create these applications.

4.1 Hub

Hub is a device which floods the packets that it receives on a specific port. As in Fig. 7, when the hub receives the packet at port 1, it floods the packets to all others ports, in this case port 2 and port 3. For all the applications, we will create the topology using this simple command

```
# mn --topo single,3 --controller=remote
```

In one terminal. On other terminal in case of hub, we will run the command

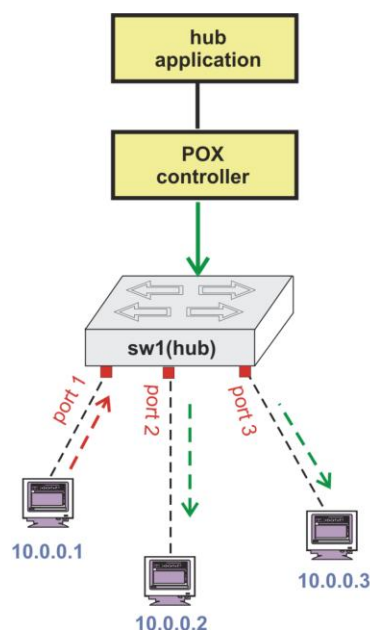
```
# ./pox.py pox.forwarding.hub
```

We capture the traffic using tcpdump command. When we are sending ping from h1 to h3, packets are received on h2 also, thus verifying the default behavior of hub.

4.2 Switch

Hub is an unintelligent device. On the other hand, switch is intelligent device as shown in Fig. 8. It maintains a mac table consisting of entries showing the relationship between ports and the mac address of the host connected to that port. As can be seen in figure 8, when we are pinging from h1 to h3, no packet is received on h2. That is we are successful in testing switch application on topology created on mininet. The command for running switch application is

```
# ./pox.py pox.forwarding.l2_learning
```

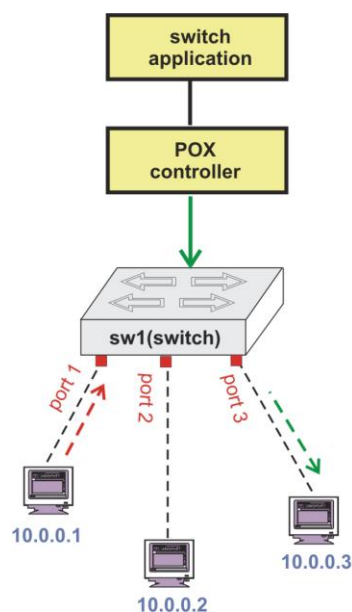


```
"Node: h1" (on mininet-vm)
root@mininet-vm:~# ping -c1 10.0.0.3
PING 10.0.0.3 (10.0.0.3) 56(84) bytes of data:
64 bytes from 10.0.0.3: icmp_seq=1 ttl=64 time=1.47 ms
--- 10.0.0.3 ping statistics ---
1 packets transmitted, 1 received, 0% packet loss, time 0ms
rtt min/avg/max/ndev = 1.477/1.477/1.477/0.000 ms
root@mininet-vm:~#

"Node: h2" (on mininet-vm)
root@mininet-vm:~# tcpdump -i h2-eth0
tcpdump: verbose output suppressed, use -v or -vv for full protocol decode
listening on h2-eth0, link-type EN10MB (Ethernet), capture size 65535 bytes
01:49:20.664417 ARP, Request who-has 10.0.0.3 tell 10.0.0.1, length 28
01:49:20.664787 ARP, Reply 10.0.0.3 is-at 00:00:00:00:00:03 (oui Ethernet), length 28
00:03 (oui Ethernet), length 28
01:49:20.664991 IP 10.0.0.1 > 10.0.0.3: ICMP echo request, id 1909, seq 1, length 64
01:49:20.665221 IP 10.0.0.3 > 10.0.0.1: ICMP echo reply, id 1909, seq 1, length 64
01:49:25.670814 ARP, Request who-has 10.0.0.1 tell 10.0.0.3, length 28

"Node: h3" (on mininet-vm)
root@mininet-vm:~# tcpdump -i h3-eth0
tcpdump: verbose output suppressed, use -v or -vv for full protocol decode
listening on h3-eth0, link-type EN10MB (Ethernet), capture size 65535 bytes
01:49:20.664417 ARP, Request who-has 10.0.0.3 tell 10.0.0.1, length 28
01:49:20.664787 ARP, Reply 10.0.0.3 is-at 00:00:00:00:00:03 (oui Ethernet), length 28
01:49:20.664991 IP 10.0.0.1 > 10.0.0.3: ICMP echo request, id 1909, seq 1, length 64
01:49:20.665221 IP 10.0.0.3 > 10.0.0.1: ICMP echo reply, id 1909, seq 1, length 64
01:49:25.670814 ARP, Request who-has 10.0.0.1 tell 10.0.0.3, length 28
01:49:25.670847 ARP, Reply 10.0.0.1 is-at 00:00:00:00:00:01 (oui Ethernet), length 28
```

Fig. 7 hub application



```
"Node: h1" (on mininet-vm)
root@mininet-vm:~# ping -c1 10.0.0.3
PING 10.0.0.3 (10.0.0.3) 56(84) bytes of data:
64 bytes from 10.0.0.3: icmp_seq=1 ttl=64 time=9.87 ms
--- 10.0.0.3 ping statistics ---
1 packets transmitted, 1 received, 0% packet loss, time 0ms
rtt min/avg/max/ndev = 9.872/9.872/9.872/0.000 ms
root@mininet-vm:~#

"Node: h2" (on mininet-vm)
root@mininet-vm:~# tcpdump -i h2-eth0
tcpdump: verbose output suppressed, use -v or -vv for full protocol decode
listening on h2-eth0, link-type EN10MB (Ethernet), capture size 65535 bytes
01:58:32.522230 IP 10.0.0.1 > 10.0.0.3: ICMP echo request, id 1923, seq 1, length 64
01:58:32.522291 IP 10.0.0.3 > 10.0.0.1: ICMP echo reply, id 1923, seq 1, length 64
01:58:37.525830 ARP, Request who-has 10.0.0.1 tell 10.0.0.3, length 28
01:58:37.561618 ARP, Request who-has 10.0.0.3 tell 10.0.0.1, length 28
01:58:37.561647 ARP, Reply 10.0.0.3 is-at 00:00:00:00:00:03 (oui Ethernet), length 28
01:58:37.564163 ARP, Reply 10.0.0.1 is-at 00:00:00:00:00:01 (oui Ethernet), length 28

"Node: h3" (on mininet-vm)
root@mininet-vm:~# tcpdump -i h3-eth0
tcpdump: verbose output suppressed, use -v or -vv for full protocol decode
listening on h3-eth0, link-type EN10MB (Ethernet), capture size 65535 bytes
01:58:32.522230 IP 10.0.0.1 > 10.0.0.3: ICMP echo request, id 1923, seq 1, length 64
01:58:32.522291 IP 10.0.0.3 > 10.0.0.1: ICMP echo reply, id 1923, seq 1, length 64
01:58:37.525830 ARP, Request who-has 10.0.0.1 tell 10.0.0.3, length 28
01:58:37.561618 ARP, Request who-has 10.0.0.3 tell 10.0.0.1, length 28
01:58:37.561647 ARP, Reply 10.0.0.3 is-at 00:00:00:00:00:03 (oui Ethernet), length 28
01:58:37.564163 ARP, Reply 10.0.0.1 is-at 00:00:00:00:00:01 (oui Ethernet), length 28
```

Fig. 8 switch application

4.3. Load Balancer

Load Balancer is an application that is used to distribute load to different servers based on different strategies such as random, round robin, least connections etc. Default load balancer included with pox controller is based on random strategy. For testing this application we created the same topology as used above using mininet. Then we ran the application in other terminal by using the command

`./pox.py pox.misc.ip_loadbalancer -ip=10.0.0.254 -servers=10.0.0.2,10.0.0.3`

As shown in Fig. 9, when we try to access web server on virtual ip address 10.0.0.254 from host h1, it first send the request to h3 and then 3 requests to h2, then again to h3 on based of random strategy. Thus mininet enables us to test different application. Same code can be used on real hardware without making any changes. Thus Mininet can be used as perfect test bed for testing SDN Applications.

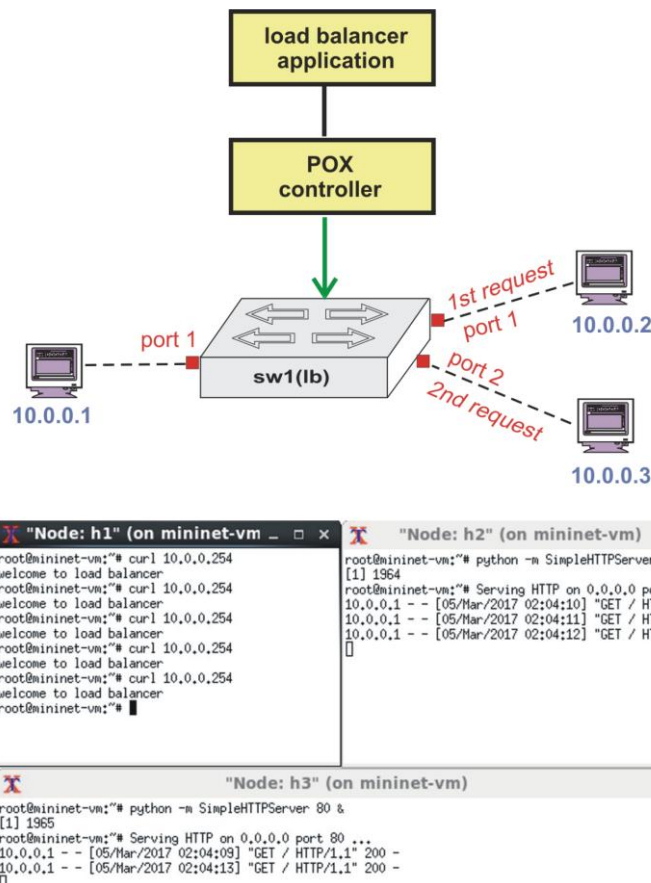


Fig. 9 load balancer application

V. MININET ADVANTAGES AND DISADVANTAGES

5.1. Advantages

- You can create simple default topologies consisting of 1 switch, 2 hosts as well as large topologies consisting of thousands of switches and hosts.
- You can also create custom topologies by writing few lines of python code.
- Mininet can be installed on physical or virtual machine. It can be installed on even raspberry pi.
- No need for any modification in network application code. Any code that will run on mininet, will run on real hardware.
- Creating networks in mininet takes very little time. Performance is good.
- Using mininet is very easy. Learning curve is very fast.

5.2. Disadvantages

- When creating very large topologies, performance suffers.

VI. CONCLUSION

Mininet can be used as a tested for doing SDN research. It enables us to create standard topologies as well as custom topologies by writing few lines of python code. It is much better than simulators where you cannot run real application code. One issue with Mininet could be performance when creating very large scale topologies. Future work could involve extending mininet code to work in clustering mode so that performance can be further increased.

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